Masked Hard-Attention Transformers and Boolean RASP Recognize Exactly the Star-Free Languages

Dana Angluin, David Chiang, Andy Yang

How's My LISP Syntax?

```
> (defun even(num) (= (mod num 2) 0))
> (filter '(6 4 3 5 2) #'even)
> (6 4 2)
```

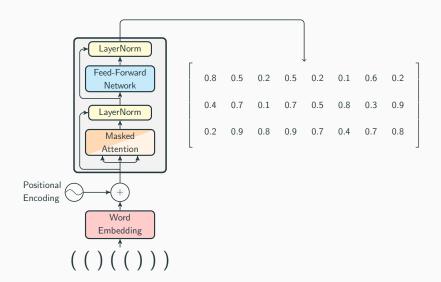
How's My LISP Syntax?

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(defun even(num)(= (mod num 2) 0))
(filter '(6 4 3 5 2) #'even)
```

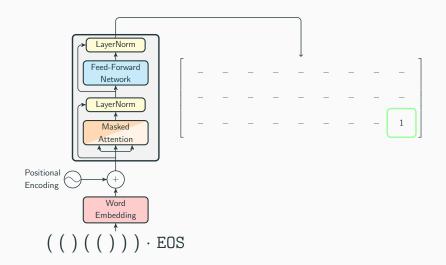
How's My LISP Syntax?

(()(()))

Transformer Encoders



Transformer Encoders as Formal Language Recognizers



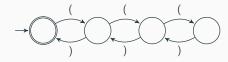
Bounded-Depth Dyck Language

Dyck-1 of depth 3

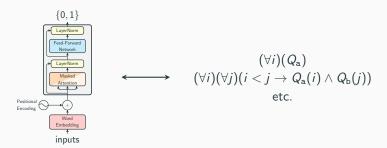
- = strings of parentheses, balanced and nested up to 3 deep
- = strings where the number of ('s is equal to the number of)'s, and every prefix contains 0–3 more ('s than)'s

Examples:

- accepted: ()() ✓
- accepted: (())() ✓
- accepted: (()(())) ✓
- rejected: (((()))) X
- rejected: ()()(**X**



The Big Picture: Expressivity and Logic



What languages are recognized by transformer encoders?

What languages are defined by logical formulas?

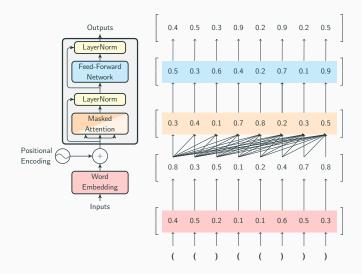
For a survey of papers in this area (including this one): Strobl et al. [4], "Transformers as Recognizers of Formal Languages: A Survey on Expressivity"

The Big Picture

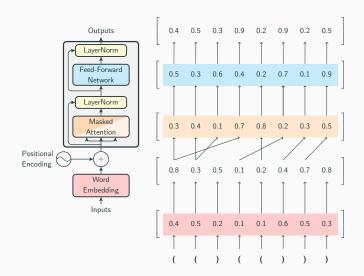
Questions to Consider

- Expressivity what can transformers do under perfect conditions?
- Learnability what can transformers learn to do in real life?
- Interpretability how can we know what transformers have learned?
- Improvements how can we augment the architecture?

Standard Attention



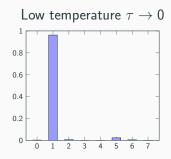
Hard Attention Simplifies Our Analysis

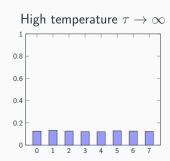


More on Hard Attention

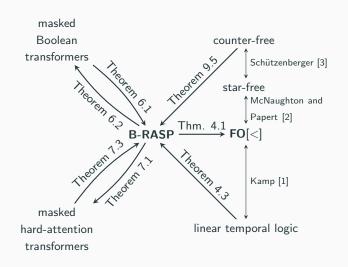
We can think of SOFTMAX as parameterized with a "temperature". At low temperature, it closely approximates ARGMAX - selecting a single input – which is much easier to manage!

$$\text{SOFTMAX}(x_i, \tau) = \frac{e^{x_i/\tau}}{\sum_{j=0}^{n} e^{x_j/\tau}}$$





Our Results



Star-Free Languages

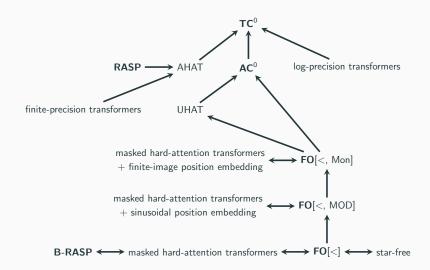
Formula	Language
Dyck-1 of Depth 3	$(),(()),((())),()(),(())(),\dots$
(ab)*	$\epsilon,$ ab, abab, ababab, \dots
$\overline{\Sigma^*aa\Sigma^*}$	$\epsilon,$ a, ab, ba, bb, abb, bab, bba, bbb,
aΣ*b	ab, aab, abb, aaab, aabb,

FO[<] and Star-Free

Using a theorem of McNaughton and Papert [2], the Star-Free languages are exactly those described by First-Order Logic

Formula	Language
$\exists x. Q_1(x)$	Strings containing a 1
$\forall x. \neg Q_0(x)$	Strings not containing any 0's
$\exists x. Q_0(x) \land \forall y.y > x$	String's starting with a 0
$\exists x. \exists y. \exists z. x < y < z$ $\land Q_0(x) \land Q_0(y) \land Q_1(z)$	String's containing the subsequence 001

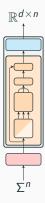
Contextualizing Our Results



Thank You

Stephen Bothwell, Darcey Riley, Ken Sible, Aarohi Srivastava, Lena Strobl, and Chihiro Taguchi!

Questions?



- With some generous assumptions, we can prove what transformers are capable of perfect conditions!
- Masked hard-attention transformer as a "base case" to build upon
- Very rich computational equivalences
- Ultimate goal: understand rigorously the capabilities and limitations of transformers

References

- [1] Johan Anthony Willem Kamp. Tense Logic and the Theory of Linear Order. PhD thesis, University of California, Los Angeles, 1968. URL https://www.proquest.com/docview/302320357.
- [2] Robert McNaughton and Seymour Papert. Counter-Free Automata. Number 65 in M.I.T. Press Research Monographs. The M.I.T. Press, 1971. ISBN 9780262130769. URL https://archive.org/embed/CounterFre_00_McNa.
- [3] M.P. Schützenberger. On finite monoids having only trivial subgroups. *Information and Control*, 8(2):190–194, 1965. DOI 10.1016/S0019-9958(65)90108-7.

References ii

[4] Lena Strobl, William Merrill, Gail Weiss, David Chiang, and Dana Angluin. Transformers as recognizers of formal languages: A survey on expressivity. *arXiv preprint arXiv:2311.00208*, 2023.